

EGC442 Introduction to Computer Architecture

Test 1

First Name: _____ **Last Name:** _____

- You are allowed to print a copy of the instruction sheet.
<http://www.engr.newpaltz.edu/~bai/EGC442/Instructin%20set%20reference.pdf>
- Your submission must be in a single PDF file
- Make sure you submit before the deadline of 1:45 PM. I will not accept late submission by email.
- You must adhere to the honor code. Any evidence of misconduct will be dealt with strictly per syllabus.

Question 1 (10 points)

Convert -12.75_{10} to single precision IEEE 754 format. Show your representation in binary.

$$\begin{array}{ccccccc} & 8 & 4 & 2 & 1 & . & 5 & . & 25 \\ 1 & 1 & 0 & 0 & . & 1 & 1 & & \\ \hline & & & & & & & & \\ \text{Exp} & = & 127 & + & 3 & = & 130 \end{array}$$

$$1.10011 \times 2^3$$

1	1000	0010	10011000	-----	0
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Question 2 (10 points)

A program runs in 100 seconds. Multiply and divide operations are responsible for 30 and 40 of those seconds, respectively. Would it be possible to run the program 2 times faster?

- 3 a. By improving the multiplication alone. If yes, by what factor?
 3 b. By improving the division alone. If yes, by what factor?
 4 c. By improving the multiplication and division by the same factor. If yes, by what factor?

$$100 = 30 + 40 + 30$$

mult
Divide
other

2X ↓

a. $50 \neq 30 + 70$
 not possible

b. $50 = 60 + 30$
 not possible

c. $50 = 20 + 30$

$$\frac{70}{20} = 3.5X$$

Question 3 (15 points)

Suppose we have two implementations of the same instruction set architecture. Computer A has a clock cycle time of 200 ps and a CPI of 2.0 a program with 2×10^6 instructions, and computer B has a clock cycle time of 350 ps and a CPI of 1.2 for the same program.

- 3 a. How many clock cycles each computer execute?
 4 b. What is the execution time of each processor?
 4 c. What is the MIPS rating of each processor?
 4 d. Which processor is a better processor? Why?

$$T_A = 200 \times 10^{-12} \quad CPI_A = 2 \quad N_A = 2 \times 10^6$$

$$T_B = 350 \times 10^{-12} \quad CPI_B = 1.2$$

a. $CLK_A = 2 \times 10^6 \times 2 = 4 \times 10^6$
 $CLK_B = 2 \times 10^6 \times 1.2 = 2.4 \times 10^6$

b. $T_{exeA} = 2 \times 10^6 \text{ instr} \times 2 \frac{\text{clk}}{\text{instr}} \times 200 \times 10^{-12} \frac{\text{sec}}{\text{clk}} = 800 \times 10^{-6} \text{ sec.}$
 $T_{exeB} = 2 \times 10^6 \text{ instr} \times 1.2 \frac{\text{clk}}{\text{instr}} \times 350 \times 10^{-12} \frac{\text{sec}}{\text{clk}} = 840 \times 10^{-6} \text{ sec.}$

c. $MIPS_A = \frac{2 \times 10^6}{10^6 \times 800 \times 10^{-6}} = 2500$
 $MIPS_B = \frac{2 \times 10^6}{10^6 \times 840 \times 10^{-6}} = 2381$

d. A is better due to lower execution time

Question 4 (15 points)

Replace the following C code by a set of equivalent MIPS instructions. Assume x, y, z, and i are in registers \$t4, \$t5, \$t6, and \$a2, respectively, and the starting address for array a is in register \$s4.

```

if (x == y)
    a[i] = a[i] + a[i];
else
    a[i+1] = a[i] - z;

```

Handwritten annotations: \$t4 ← if (x == y), \$t5 ← a[i] = a[i] + a[i];, \$s4 ← else, \$t6 ← a[i+1] = a[i] - z;

```

sll    $a2, $a2, 2
add    $s4, $s4, $a2
lw     $t0, 0($s4)
beq    $t4, $t5, ELSE
sub    $t0, $t0, $t6
sw     $t0, 4($s4)
j      out
ELSE:  add    $t0, $t0, $t0
       sw     $t0, 0($s4)

```

out:

0-5 Small mistake
5-10 Med.
10-15 Big mistakes

Question 5 (10 points)

Assume \$s3 has 5002, and words addressed 5000..5002 have the data shown:

5000: 0x99

5001: 0x77

5002: 0x23

5003: 0x23

5004: 0x6E

5005: 0x34

5006: 0x13

5007: 0x34

a. What address will be computed by lw \$t0, -2(\$s3)

b. What value will be put in \$t1 by lw \$t1, 2(\$s3)

0x5000

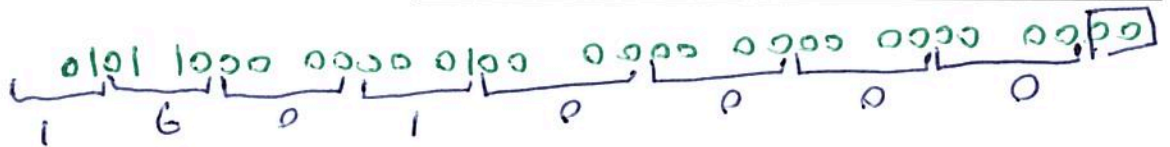
5004

\$t1 = 0x6E341334

Question 6 (15 points)

The following code segment is partially converted to machine code. Complete the conversion by placing appropriate values in the empty machine code fields of *sub*, *beq*, and *j* instructions.

<i>Loop: sub \$t1, \$t1, \$a2</i>	0x58040000	0	9	6	9	0	0x22
<i>lw \$t0, 0(\$t1)</i>	0x58040004	35	9	8	0		
<i>beq \$t0, \$s5, Exit</i>	0x58040008	4	8	21	2	(3)	
<i>addi \$s5, \$s5, 1</i>	0x5804000C	8	21	21	1		
<i>j Loop</i>	0x58040010	2	0x16010000				(3)
<i>Exit: ...</i>	0x58040014	xxxxxx					

**Question 7 (10 points)**

Compare the number of gate delays for the critical paths of the following 32-bit adders

- 2 Only using ripple carry
- 4 Using carry lookahead at the first level and ripple carry at between them
- 4 Using carry lookahead at the first and second level and ripple carry between second levels.

i. $32 \times 2 = 64$ gate delays

ii. $32/4 = 8$ modules

$8 \times 3 = 24$ gate delays

iii.

$$\frac{32}{16} = 2$$

$2 \times 5 = 10$ gate delays

