EGC442
Problem Set 9
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1) Refer to the refined multiplication hardware figure below

a) The refined multiplication hardware halves the width of the Multiplicand register from 64-bits to 32-bits.
© True
O False
b) The Multiplier register is removed and placed inside of the $\qquad$ register.
$\bigcirc$ Product
O Multiplicand
c) The ALU adds the 64-bit Product and 32-bit Multiplicand, and then stores the result into the Product register.
O True
O False
2) Consider the multiplication of $510 \times 12_{10}$, or $0101_{2} \times 1100_{2}$. Fill in the missing values for each of the steps labeled according to COD Figure 3.4 (The first multiplication algorithm ...). A copy of the multiplication algorithm figure is shown below to the right.

| Iteration | Step | Multiplier | Multiplicand | Product |
| :--- | :--- | :---: | :---: | :--- |
| 0 | Initial values | 1100 | 00000101 | 00000000 |
|  | 1: $0 \Rightarrow$ No operation | 1100 | 00000101 | 00000000 |
|  | 2: Shift left Multiplicand | 1100 | 00001010 | 00000000 |
|  | 3: Shift right Multiplier | 0110 | 00001010 | 00000000 |
| 3 | (a) | 0110 | 00001010 | 00000000 |
|  | 2: Shift left Multiplicand | 0110 | (b) | 00000000 |
|  | 3: Shift right Multiplier | (c) |  | 00000000 |
| 4 | (d) |  |  | 00010100 |
|  | 2: Shift left Multiplicand |  | 00101000 | 00010100 |
|  | 3: Shift right Multiplier | 0001 | 00101000 | 00010100 |
|  | 1a: 1 $\Rightarrow$ Prod = Prod + Mcand | 0001 | 00101000 | (e) |
|  | 2: Shift left Multiplicand | 0001 | 01010000 |  |
|  | 3: Shift right Multiplier | 0000 | 01010000 |  |



- 00010100
- 1a: 1 ==> Prod = Prod + Mcand
- 00111100
- 0011

1: $0==>$ No operation
(a)
(b)
(c)
(d)
(e)
3)
a. The multiplication hardware supports signed multiplication.

O True
O False
b. The 32-bit registers, called Hi and Lo, combine to form a 64-bit product register.

O True
$\bigcirc$ False
c. The multiply (mult) instruction ignores overflow, while the multiply unsigned (multu) instruction detects overflow.
O True
O False
4)
a. A calculation that leads to a number being too large to represent is called $\qquad$ .
overflow
O underflow
O a fraction
b. Increasing the size of the $\qquad$ used to represent a floating-point number impacts the number's precision.
fraction
O exponent
d. A $\qquad$ precision floating-point number is represented with two MIPS words.
O single
O double
5. Show the IEEE 754 binary representation of the number +0.375 ten in single precision:

```
- 112 / 2 or 0.011 2
- }12
- . }10000000000000000000000
3 / }8\mathrm{ or 3/2
1.1 (two }\times\mp@subsup{2}{}{-2
    0
- (-1)}\mp@subsup{)}{}{0}\times(1+.10000000000000000000 0000) \times 2(125-127)
```

Rewrite as a fraction
Rewrite as a binary number

Rewrite as normalized scientific notation
$\mathrm{S}=$ ?
Exponent = ?
Fraction $=$ ?
IEEE 754 binary single precision representation
6. Show the IEEE 754 binary representation of the number -0.9375 ten in double precision:

```
-1
    1022
- (-1)}\mp@subsup{)}{}{1}\times(1+.11100000\ldots0000)\times\mp@subsup{2}{}{(1022-1023)
- 1111 two / 24 or 0.1111 two
- . }11100000 ... 0000
- 15/16 or 15/24
1.111 two }\times\mp@subsup{2}{}{-1
```

Rewrite as a fraction
Rewrite as a binary number
Rewrite as normalized scientific notation
$\mathrm{S}=$ ?
Exponent $=$ ?
Fraction $=$ ?
IEEE 754 binary double precision representation
7. Convert the single precision binary floating-point representation to decimal.

| 31 | 30 | 29 | 28 | 27 | 26 | 25 | 24 | 23 | 22 | 21 | 20 | 19 | 18 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

8. Add the following numbers using the floating-point addition algorithm. Assume 4 bits of precision.
a. $1.010 \times 2^{-3}+0.011 \times 2^{-3}=$ ?

- 1.101
© $1.101 \times 2^{-6}$
© $1.101 \times 2^{-3}$
b. $1.001 \times 2^{-4}+1.000 \times 2^{-6}=$ ?
o $10.001 \times 2^{-4}$
© $1.011 \times 2^{-4}$
c. $1.000 \times 2^{3}+0.011 \times 2^{5}=$ ?

C $1.010 \times 2^{4}$
C $0.101 \times 2^{5}$
C $10.001 \times 2^{5}$
9. Multiply $-14_{\text {ten }}$ and $-0.25_{\text {ten }}$, or $-1.110 \times 2^{3} \times-1.000 \times 2^{-2}$. Assume 4 bits of precision.

```
-1.110000
-1.110000}\mp@subsup{0}{\mathrm{ two }}{\times}\times\mp@subsup{2}{}{1
- 1.1100 two * 2 }\mp@subsup{}{}{1
- 3+(-2)=1
- 2
- 3.5ten
```

Adding the non-biased exponents of the operands
Multiply the significands:
$1.110 \times 1.000=$ ?
Product $=1.110000 \times$ ?
Normalize the product
Round the product
Set the sign of the product: ? $1.1100_{\text {two }} \times 2^{1}$
$-14_{\text {ten }} \times-0.25_{\text {ten }}=$ ?

